

Presenting isomorphic finitely presented modules by equivalent matrices: a constructive approach

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Linear systems and finitely presented left D -modules

- ◇ D ring of functional operators, $R \in D^{q \times p}$, and \mathcal{F} a left D -module
- ◇ Consider the linear system $\ker_{\mathcal{F}}(R.) = \{\eta \in \mathcal{F}^p \mid R\eta = 0\}$.
- ◇ We associate the following finitely presented left D -module:

$$M = D^{1 \times p} / (D^{1 \times q} R),$$

given by the finite presentation:

$$\begin{array}{ccccccc} D^{1 \times q} & & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\pi} & M & \longrightarrow 0, \\ \lambda = (\lambda_1, \dots, \lambda_q) & \longmapsto & & \lambda R & & & \end{array}$$

- ◇ Malgrange's isomorphism: $\ker_{\mathcal{F}}(R.) \cong \text{hom}_D(M, \mathcal{F})$
- ◇ **Algebraic analysis**: the linear system $\ker_{\mathcal{F}}(R.)$ can be studied by means of the finitely presented left D -module M .

Example in the theory of linear elasticity (*Pommaret'01*)

$$\diamond (S) \begin{cases} \partial_1 \xi_1 = 0, \\ \frac{1}{2} (\partial_2 \xi_1 + \partial_1 \xi_2) = 0, \\ \partial_2 \xi_2 = 0. \end{cases} \quad (S') \begin{cases} \partial_1 \zeta_1 = 0, & \partial_2 \zeta_1 - \zeta_2 = 0, \\ \partial_1 \zeta_2 = 0, & \partial_1 \zeta_3 + \zeta_2 = 0, \\ \partial_2 \zeta_3 = 0, & \partial_2 \zeta_2 = 0. \end{cases}$$

◊ We consider the ring $D = \mathbb{Q}[\partial_1, \partial_2]$ and we have

$$(S) \Leftrightarrow \overbrace{\begin{pmatrix} \partial_1 & 0 \\ \frac{1}{2} \partial_2 & \frac{1}{2} \partial_1 \\ 0 & \partial_2 \end{pmatrix}}^R \begin{pmatrix} \xi_1 \\ \xi_2 \end{pmatrix} = 0, \quad (S') \Leftrightarrow \overbrace{\begin{pmatrix} \partial_1 & 0 & 0 \\ \partial_2 & -1 & 0 \\ 0 & \partial_1 & 0 \\ 0 & 1 & \partial_1 \\ 0 & 0 & \partial_2 \\ 0 & \partial_2 & 0 \end{pmatrix}}^{R'} \begin{pmatrix} \zeta_1 \\ \zeta_2 \\ \zeta_3 \end{pmatrix} = 0.$$

◊ We associate $M = D^{1 \times 2} / (D^{1 \times 3} R)$ and $M' = D^{1 \times 3} / (D^{1 \times 6} R')$.

Two theorems about isomorphisms and equivalences

Theorem (Fitting, 1936)

Two matrices $R \in D^{q \times p}$ and $R' \in D^{q' \times p'}$ presenting isomorphic left D -modules can be inflated with blocks of 0 and I to get equivalent matrices presenting the same left D -modules.

Theorem (Warfield, 1978)

If two positive integers s and r satisfy

$$\begin{aligned} s &\leq \min(p + q', q + p'), & \text{sr}(D) &\leq \max(p + q' - s, q + p' - s), \\ r &\leq \min(p, p'), & \text{sr}(D) &\leq \max(p - r, p' - r), \end{aligned}$$

then we can remove s blocks of zeros and r blocks of identity.

◇ Goal of the talk: give constructive versions of both theorems

Example in the theory of linear elasticity

- ◊ We will prove that $M = D^{1 \times 2} / (D^{1 \times 3} R) \cong M' = D^{1 \times 3} / (D^{1 \times 6} R')$.
- ◊ Constructive version of Fitting's theorem: compute $X \in GL_5(D)$ and $Y \in GL_{14}(D)$ such that:

$$\left(\begin{array}{cc|ccc}
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 \hline
 1 & 0 & 0 & 0 & 0 \\
 0 & 1 & 0 & 0 & 0 \\
 \hline
 0 & 0 & \partial_1 & 0 & 0 \\
 0 & 0 & \partial_2 & -1 & 0 \\
 0 & 0 & 0 & \partial_1 & 0 \\
 0 & 0 & 0 & 1 & \partial_1 \\
 0 & 0 & 0 & 0 & \partial_2 \\
 0 & 0 & 0 & \partial_2 & 0
 \end{array} \right) = Y^{-1} \left(\begin{array}{cc|ccc}
 \partial_1 & 0 & 0 & 0 & 0 \\
 \frac{1}{2} \partial_2 & \frac{1}{2} \partial_1 & 0 & 0 & 0 \\
 0 & \partial_2 & 0 & 0 & 0 \\
 \hline
 0 & 0 & 1 & 0 & 0 \\
 0 & 0 & 0 & 1 & 0 \\
 0 & 0 & 0 & 0 & 1 \\
 \hline
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0
 \end{array} \right) X.$$

Example in the theory of linear elasticity

◇ Constructive version of Warfield's theorem (slight generalization):
compute $X' \in GL_3(D)$ and $Y' \in GL_7(D)$ such that:

$$\begin{pmatrix} 0 & 0 & 0 \\ \partial_1 & 0 & 0 \\ \partial_2 & -1 & 0 \\ 0 & \partial_1 & 0 \\ 0 & 1 & \partial_1 \\ 0 & 0 & \partial_2 \\ 0 & \partial_2 & 0 \end{pmatrix} = Y'^{-1} \begin{pmatrix} \partial_1 & 0 & | & 0 \\ \frac{1}{2} \partial_2 & \frac{1}{2} \partial_1 & | & 0 \\ 0 & \partial_2 & | & 0 \\ \hline 0 & 0 & | & 1 \\ 0 & 0 & | & 0 \\ 0 & 0 & | & 0 \\ 0 & 0 & | & 0 \end{pmatrix} X'.$$

Part I

Definition of isomorphic finitely presented left D -modules in terms of matrix equalities

Homomorphisms in terms of matrix equality

- ◇ Let D be a ring of functional operators.
- ◇ Let $R \in D^{q \times p}$, $R' \in D^{q' \times p'}$ be two matrices.
- ◇ We have the following **commutative exact diagram**:

$$\begin{array}{ccccccc} D^{1 \times q} & \xrightarrow{\cdot R} & D^{1 \times p} & \xrightarrow{\pi} & M & \longrightarrow & 0 \\ \downarrow \cdot Q & & \downarrow \cdot P & & \downarrow f & & \\ D^{1 \times q'} & \xrightarrow{\cdot R'} & D^{1 \times p'} & \xrightarrow{\pi'} & M' & \longrightarrow & 0. \end{array}$$

$\exists f : M \rightarrow M' \iff \exists P \in D^{p \times p'}, Q \in D^{q \times q'}$ such that:

$$R P = Q R'.$$

Moreover, we have $f(\pi(\lambda)) = \pi'(\lambda P)$, for all $\lambda \in D^{1 \times p}$.

Isomorphisms in terms of matrix equalities

◇ $f \in \text{hom}_D(M, M')$ given by P and Q such that $RP = QR'$.

◇ $f \in \text{iso}_D(M, M')$ if and only if there exist $P' \in D^{p' \times p}$, $Q' \in D^{q' \times q}$, $Z \in D^{p \times q}$ and $Z' \in D^{p' \times q'}$ satisfying the following relations:

$$R'P' = Q'R, \quad PP' + ZR = I_p, \quad P'P + Z'R' = I_{p'}.$$

Then, there exist $Z_2 \in D^{q \times r}$ and $Z'_2 \in D^{q' \times r'}$ such that:

$$QQ' + RZ + Z_2R_2 = I_q, \quad Q'Q + R'Z' + Z'_2R'_2 = I_{q'},$$

where $\ker_D(.R) = D^{1 \times r} R_2$ and $\ker_D(.R') = D^{1 \times r'} R'_2$.

Implementations

- ◇ **All the matrices** appearing in the relations defining (iso)morphisms **can be computed from R and R'** using (non-commutative) Gröbner basis calculations.
- ◇ Implementations:
 - 1 **Maple package OREMORPHISMS** (C.Q.) based on OREMODULES (*Chyzak, Q., Robertz*);
 - 2 **Mathematica package OREALGEBRAICANALYSIS** (*C., Q., Tönso*) based on HOLONOMICFUNCTIONS (*Koutschan*);
 - 3 **CapAndHoma1g** (*Barakat et al.*).

Example in the theory of linear elasticity

◇ We have $RP = QR'$ with:

$$P = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}, \quad Q = \frac{1}{2} \begin{pmatrix} 2 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 2 & 0 \end{pmatrix}.$$

◇ We have $R'P' = Q'R$, $PP' + ZR = I_p$, and $P'P + Z'R' = I_{p'}$ with:

$$P' = \begin{pmatrix} 1 & 0 \\ \partial_2 & 0 \\ 0 & 1 \end{pmatrix}, \quad Q' = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \\ \partial_2 & 0 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 1 \\ 0 & 2\partial_2 & -\partial_1 \end{pmatrix},$$
$$Z = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}, \quad Z' = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}.$$

◇ This proves that $M = D^{1 \times 2} / (D^{1 \times 3} R) \cong M' = D^{1 \times 3} / (D^{1 \times 6} R')$.

Part II

Constructive version of Fitting's theorem

Statement of the problem

- ◇ Let $M = D^{1 \times p} / (D^{1 \times q} R)$ and $M' = D^{1 \times p'} / (D^{1 \times q'} R')$.
- ◇ Assume that $M \cong M'$

Theorem (Fitting, 1936)

Two matrices $R \in D^{q \times p}$ and $R' \in D^{q' \times p'}$ presenting isomorphic left D -modules can be inflated with blocks of 0 and I to get equivalent matrices presenting the same left D -modules.

- ◇ Goal: Compute inflations L and L' of R and R' with blocks of 0 and I and two unimodular matrices X and Y such that:
 - L and L' respectively define a finite presentation of M and M' ;
 - L and L' are equivalent matrices, i.e., we have $L' = Y^{-1} L X$.
- ◇ We give explicit formulas for X and Y in terms of the matrices defining the isomorphism.

Explicit Fitting's theorem

◇ Let $\begin{cases} n = q + p' + p + q' \\ m = p + p' \end{cases}$ and define the matrices:

1

$$X = \begin{pmatrix} I_p & P \\ -P' & I_{p'} - P' P \end{pmatrix} \in \mathrm{GL}_m(D), \quad X^{-1} = \begin{pmatrix} I_p - P P' & -P \\ P' & I_{p'} \end{pmatrix}.$$

2

$$Y = \begin{pmatrix} I_q & 0 & R & Q \\ 0 & I_{p'} & -P' & Z' \\ -Z & P & 0 & P Z' - Z Q \\ -Q' & -R' & 0 & Z'_2 R'_2 \end{pmatrix} \in \mathrm{GL}_n(D),$$

with inverse given by

$$Y^{-1} = \begin{pmatrix} Z_2 R_2 & 0 & -R & -Q \\ P' Z - Z' Q' & 0 & P' & -Z' \\ Z & -P & I_p & 0 \\ Q' & R' & 0 & I_{q'} \end{pmatrix}.$$

Explicit Fitting's theorem

$$\diamond \text{ Let } L = \begin{pmatrix} R & 0 \\ 0 & I_{p'} \\ 0 & 0 \\ 0 & 0 \end{pmatrix} \in D^{n \times m}, \quad L' = \begin{pmatrix} 0 & 0 \\ 0 & 0 \\ I_p & 0 \\ 0 & R' \end{pmatrix} \in D^{n \times m}.$$

\diamond The following commutative exact diagram holds:

$$\begin{array}{ccccc} 0 & & 0 & & 0 \\ \downarrow & & \downarrow & & \downarrow \\ D^{1 \times n} & \xrightarrow{\cdot L} & D^{1 \times m} & \xrightarrow{\pi \oplus 0_{p'}} & M \rightarrow 0 \\ \downarrow \cdot Y & & \downarrow \cdot X & & \downarrow f \\ D^{1 \times n} & \xrightarrow{\cdot L'} & D^{1 \times m} & \xrightarrow{0_p \oplus \pi'} & M' \rightarrow 0 \\ \downarrow & & \downarrow & & \downarrow \\ 0 & & 0 & & 0 \end{array}$$

and L and L' are equivalent matrices, i.e.,

$$\begin{pmatrix} 0 & 0 \\ 0 & 0 \\ I_p & 0 \\ 0 & R' \end{pmatrix} = Y^{-1} \begin{pmatrix} R & 0 \\ 0 & I_{p'} \\ 0 & 0 \\ 0 & 0 \end{pmatrix} X.$$

Example in the theory of 2D linear elasticity

◇ We get

$$\left(\begin{array}{cc|ccc} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ \hline 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ \hline 0 & 0 & \partial_1 & 0 & 0 \\ 0 & 0 & \partial_2 & -1 & 0 \\ 0 & 0 & 0 & \partial_1 & 0 \\ 0 & 0 & 0 & 1 & \partial_1 \\ 0 & 0 & 0 & 0 & \partial_2 \\ 0 & 0 & 0 & \partial_2 & 0 \end{array} \right) = Y^{-1} \left(\begin{array}{cc|ccc} \partial_1 & 0 & 0 & 0 & 0 \\ \frac{1}{2} \partial_2 & \frac{1}{2} \partial_1 & 0 & 0 & 0 \\ 0 & \partial_2 & 0 & 0 & 0 \\ \hline 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ \hline 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{array} \right) X.$$

Example in the theory of 2D linear elasticity

◇ $X \in GL_5(D)$ and $Y \in GL_{14}(D)$ are given by:

$$X = \begin{pmatrix} 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ -1 & 0 & 0 & 0 & 0 \\ -\partial_2 & 0 & -\partial_2 & 1 & 0 \\ 0 & -1 & 0 & 0 & 0 \end{pmatrix},$$

$$Y = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & \partial_1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & \frac{1}{2}\partial_2 & \frac{1}{2}\partial_1 & 0 & \frac{1}{2} & 0 & \frac{1}{2} & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & \partial_2 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & -1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & -\partial_2 & 0 & 0 & -1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & -1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & 0 & -\partial_1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & -\partial_2 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ -\partial_2 & 0 & 0 & 0 & -\partial_1 & 0 & 0 & 0 & -\partial_2 & \partial_1 & 1 & 0 & 0 & 0 & 0 \\ 0 & -2 & 0 & 0 & -1 & -\partial_1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & -1 & 0 & 0 & -\partial_2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & -2\partial_2 & \partial_1 & 0 & -\partial_2 & 0 & 0 & 0 & 0 & 0 & 0 & -\partial_2 & \partial_1 & 1 & 0 \end{pmatrix}.$$

Part III

Constructive version of Warfield's theorem

Statement of the problem

- Let $M = D^{1 \times p} / (D^{1 \times q} R) \cong M' = D^{1 \times p'} / (D^{1 \times q'} R')$.
- Assume that $L' = Y^{-1} L X$ as before (Fitting's theorem).

Theorem (Warfield, 1978)

If two positive integers s and r satisfy

$$\begin{aligned} s &\leq \min(p + q', q + p'), & \text{sr}(D) &\leq \max(p + q' - s, q + p' - s), \\ r &\leq \min(p, p'), & \text{sr}(D) &\leq \max(p - r, p' - r), \end{aligned}$$

then we can remove s blocks of zeros and r blocks of identity.

- Goal: Compute $X_r \in \text{GL}_{m-r}(D)$ and $Y_{s,r} \in \text{GL}_{n-s-r}(D)$ such that:

$$\begin{pmatrix} 0 & 0 \\ I_{p-r} & 0 \\ 0 & R' \end{pmatrix} = Y_{s,r}^{-1} \begin{pmatrix} R & 0 \\ 0 & I_{p'-r} \\ 0 & 0 \end{pmatrix} X_r.$$

Stable rank of a ring D

Definitions (e.g., McConnell & Robson)

- 1 $u \in D^n$ is **unimodular** if $\exists v \in D^{1 \times n}$ such that $v u = 1$.
- 2 $u = (u_1, \dots, u_n)^T \in D^n$ is **stable** if $\exists d_1, \dots, d_{n-1} \in D$ such that $(u_1 + d_1 u_n, \dots, u_{n-1} + d_{n-1} u_n)^T \in D^{n-1}$ is unimodular.
- 3 An integer r is said to be **in the stable range of D** if $\forall n > r$, a unimodular vector $u \in D^n$ is stable.
- 4 The **stable rank $\text{sr}(D)$** of D is the smallest positive integer that is in the stable range of D . If no such integer exists, then $\text{sr}(D) = +\infty$.

◇ Examples:

- 1 If D is a principal domain, then $\text{sr}(D) \geq 2$;
- 2 $\forall n \geq 1$, we have $\text{sr}(\mathbb{Q}[x_1, \dots, x_n]) = n + 1$ (*Vasershtein'71*);
- 3 $\text{sr}(A_n(k)) = 2$ and $\text{sr}(B_n(k)) = 2$ (Stafford's theorem - *Stafford'78*).

A key result based on $\text{sr}(D)$

Lemma

Let D be a ring and n, m two integers such that $\text{sr}(D) \leq m$.

Let $u \in D^{n+m+1}$ be a *unimodular* column vector such that we have:

$$v u = (v_n \quad v_{m+1}) \begin{pmatrix} u_n \\ u_m \\ u_1 \end{pmatrix} = 1,$$

where $v_n \in D^{1 \times n}$, $v_{m+1} \in D^{1 \times (m+1)}$, $u_n \in D^n$, $u_m \in D^m$, $u_1 \in D$.

Then, *there exist* $c_1 \in D$, $\tilde{u}_m \in D^m$, $\tilde{v}_m \in D^{1 \times m}$ such that we have:

$$(c_1 \quad v_n \quad \tilde{v}_m) \begin{pmatrix} u_n \\ u_m + \tilde{u}_m u_1 \end{pmatrix} = 1.$$

- ◇ *No general algorithm* for computing c_1 , \tilde{u}_m , and \tilde{v}_m in any ring D .
- ◇ But *algorithms and heuristics* are implemented for some *particular rings* D (as for instance Weyl algebras) and *allow to treat examples*.

An iterative process

- ◇ Starting point: $X \in \text{GL}_m(D)$ and $Y \in \text{GL}_n(D)$ such that $L' = Y^{-1} L X$.
- ◇ We assume that s and r satisfy the hypothesis of Warfield's theorem.
- ◇ Remove s zeros rows: From $Y_0 := Y$, we compute recursively matrices Y_1, \dots, Y_s such that: $\forall i = 1, \dots, s$,

$$Y_i \in \text{GL}_{n-i}(D), \quad \underbrace{\begin{pmatrix} 0 & 0 \\ I_p & 0 \\ 0 & R' \end{pmatrix}}_{L'_i \in D^{(n-i) \times m}} = Y_i^{-1} \underbrace{\begin{pmatrix} R & 0 \\ 0 & I_{p'} \\ 0 & 0 \end{pmatrix}}_{L_i \in D^{(n-i) \times m}} X.$$

- ◇ Remove r identity blocks: From $Y_{s,0} := Y_s$ and $X_0 := X$, we compute recursively $Y_{s,1}, \dots, Y_{s,r}$ and X_1, \dots, X_r such that: $\forall j = 1, \dots, r$,

$$Y_{s,j} \in \text{GL}_{n-s-j}(D), \quad X_j \in \text{GL}_{m-j}(D), \quad \underbrace{\begin{pmatrix} 0 & 0 \\ I_{p-j} & 0 \\ 0 & R' \end{pmatrix}}_{L'_{s,j} \in D^{(n-s-j) \times (m-j)}} = Y_{s,j}^{-1} \underbrace{\begin{pmatrix} R & 0 \\ 0 & I_{p'-j} \\ 0 & 0 \end{pmatrix}}_{L_{s,j} \in D^{(n-s-j) \times (m-j)}} X_j.$$

Procedure to remove zero rows: general description

◇ Hyp.: we have computed $Y_{i-1} \in \text{GL}_{n-i+1}(D)$ s.t. $L'_{i-1} = Y_{i-1}^{-1} L_{i-1} X$.

$$\begin{array}{ccc}
 D^{1 \times (n-i)} & \xrightarrow{\cdot L'_i} & D^{1 \times m} \\
 \uparrow \cdot G'_i & & \uparrow \cdot J_m \\
 D^{1 \times (n-i+1)} & \xrightarrow{\cdot L'_{i-1}} & D^{1 \times m} \\
 \uparrow \cdot Y_{i-1} & & \uparrow \cdot X \\
 D^{1 \times (n-i+1)} & \xrightarrow{\cdot L_{i-1}} & D^{1 \times m} \\
 \uparrow \cdot W_i & & \uparrow \cdot J_m \\
 D^{1 \times (n-i+1)} & \xrightarrow{\cdot L_{i-1}} & D^{1 \times m} \\
 \uparrow \cdot H_i & & \uparrow \cdot J_m \\
 D^{1 \times (n-i)} & \xrightarrow{\cdot L_i} & D^{1 \times m}
 \end{array}$$

$$\implies L_i X = \underbrace{(H_i W_i Y_{i-1} G'_i)}_{Y_i \in \text{GL}_{n-i}(D)} L_{i-1}, \quad Y_i^{-1} = H'_i Y_{i-1}^{-1} W_i^{-1} G_i.$$

Procedure to remove zero rows: matrix computations

- ◊ We decompose Y_{i-1} and its inverse Y_{i-1}^{-1} by blocks as follows:

$$Y_{i-1} = \begin{pmatrix} Y_{11} \\ Y_{21} \\ Y_{31} \end{pmatrix} \begin{matrix} \leftarrow q + p' \\ \leftarrow p + q' - i \\ \leftarrow 1 \end{matrix} \qquad Y_{i-1}^{-1} = \begin{pmatrix} q + p' & & 1 \\ \downarrow & & \downarrow \\ Y'_{11} & Y'_{12} & Y'_{13} \\ & \uparrow & \\ & p + q' - i & \end{pmatrix}$$

- ◊ As $\text{sr}(D) \leq p + q' - s$, the key lemma implies that **there exist** $c \in D$, $u \in D^{p+q'-i}$, $v \in D^{1 \times (p+q'-i)}$ such that, if $k = q + p' - (i - 1)$, then

$$(c(Y'_{11})_k \quad v) \begin{pmatrix} (Y_{11})_{.k} \\ (Y_{21})_{.k} + u(Y_{31})_k \end{pmatrix} = 1.$$

- ◊ Assume $c \in D$, $u \in D^{p+q'-i}$, $v \in D^{1 \times (p+q'-i)}$ have been computed.

Procedure to remove zero rows: matrix computations

◇ $W_i = \begin{pmatrix} I_{q+p'} & 0 & 0 \\ 0 & I_{p+q'-i} & u \\ 0 & 0 & 1 \end{pmatrix} \in \text{GL}_{n-i+1}(D) \rightsquigarrow$ commuting square diagram:

$$\begin{array}{ccc} D^{1 \times (n-i+1)} & \xrightarrow{\cdot L_{i-1}} & D^{1 \times m} \\ \cdot W_i^{-1} \downarrow \uparrow \cdot W_i & & \cdot l_m \downarrow \uparrow \cdot l_m \\ D^{1 \times (n-i+1)} & \xrightarrow{\cdot L_{i-1}} & D^{1 \times m} \end{array}$$

◇ Let $\tilde{l}_i = (c(Y'_{11})_k \quad v)$, $l_i = (c(Y'_{11})_k \quad v \quad 0)$, $F_i = \begin{pmatrix} Y_{11} \\ Y_{21} + u Y_{31} \end{pmatrix}$.

Procedure to remove zero rows: matrix computations

- ◇ $G_i = \begin{pmatrix} I_{n-i} - (F_i)_{.k} \tilde{\ell}_i \\ \tilde{\ell}_i \end{pmatrix}$, $H_i = (I_{n-i} - (F_i)_{.k} \tilde{\ell}_i \quad (F_i)_{.k})$ satisfy $H_i G_i = I_{n-i}$ and the following square diagrams commutes:

$$\begin{array}{ccc} D^{1 \times (n-i+1)} & \xrightarrow{\cdot L_{i-1}} & D^{1 \times m} \\ \uparrow \cdot H_i & & \uparrow \cdot J_m \\ D^{1 \times (n-i)} & \xrightarrow{\cdot L_i} & D^{1 \times m} \end{array}$$

- ◇ $G'_i = (I_{n-i+1} - (f_k^{n-i+1})^T \ell_i W_i Y_{i-1}) \begin{pmatrix} I_{k-1} & 0 \\ 0 & 0 \\ 0 & I_{p+q'} \end{pmatrix}$, $H'_i = \begin{pmatrix} I_{k-1} & 0 & 0 \\ 0 & 0 & I_{p+q'} \end{pmatrix}$, satisfy $H'_i G'_i = I_{n-i}$ and the following square diagrams commutes:

$$\begin{array}{ccc} D^{1 \times (n-i)} & \xrightarrow{\cdot L'_i} & D^{1 \times m} \\ \uparrow \cdot G'_i & & \uparrow \cdot J_m \\ D^{1 \times (n-i+1)} & \xrightarrow{\cdot L'_{i-1}} & D^{1 \times m} \end{array}$$

Procedure to remove zero rows: matrix computations

- ◇ $Y_i = H_i W_i Y_{i-1} G_i' \in GL_{n-i}(D)$ with inverse $Y_i^{-1} = H_i' Y_{i-1}^{-1} W_i^{-1} G_i$.
- ◇ The following commutative exact diagram holds

$$\begin{array}{ccccc}
 0 & & 0 & & 0 \\
 \downarrow & & \downarrow & & \downarrow \\
 D^{1 \times (n-i)} & \xrightarrow{\cdot L_i} & D^{1 \times m} & \xrightarrow{\pi \oplus 0_{p'}} & M \longrightarrow 0 \\
 \downarrow \cdot Y_i & & \downarrow \cdot X & & \downarrow f \\
 D^{1 \times (n-i)} & \xrightarrow{\cdot L_i'} & D^{1 \times m} & \xrightarrow{0_p \oplus \pi'} & M' \longrightarrow 0 \\
 \downarrow & & \downarrow & & \downarrow \\
 0 & & 0 & & 0
 \end{array}$$

and L_i and L_i' are equivalent matrices.

- ◇ Matrices given explicitly in terms of those of the previous equivalence.

Problem reduced to computing $c \in D$, $u \in D^{p+q'-i}$, $v \in D^{1 \times (p+q'-i)}$ s.t.:

$$(c(Y'_{11})_k \quad v) \begin{pmatrix} (Y_{11})_{\cdot k} \\ (Y_{21})_{\cdot k} + u(Y_{31})_k \end{pmatrix} = 1.$$

Example in the theory of 2D linear elasticity

◇ We have $q = 3$, $p = 2$, $q' = 6$, and $p' = 3$ so that $n = 14$ and $m = 5$.

◇ We have $q + p' = 6 \leq p + q' = 8$, and $\text{sr}(\mathbb{Q}[\partial_1, \partial_2]) = 3$.

◇ The integer s has to satisfy:

$$s \leq \min(p + q', q + p') = 6, \quad \text{sr}(D) \leq \max(p + q' - s, q + p' - s) = 8 - s,$$

\implies We can remove $s = 5$ zero rows.

Example in the theory of 2D linear elasticity

◇ To remove the first zero row, we are reduced to solving:

$$(c \ (0 \ 0 \ 0 \ 0 \ 0 \ 0) \ v) \left(\begin{array}{c} \left(\begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \end{array} \right) \\ \left(\begin{array}{c} 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ -\partial_1 \\ -\partial_2 \end{array} \right) + 0 \ u \end{array} \right) = 1,$$

◇ A solution is given by $c = 0$, $v = (0 \ 1 \ 0 \ 0 \ 0 \ 0)$, and $u = (0 \ 0 \ 0 \ 0 \ 0 \ 0)^T$

Example in the theory of 2D linear elasticity

- From the above formulas, we get $Y_1 \in GL_{13}(D)$ s.t. $L'_1 = Y_1^{-1} L'_1 X$.
- Applying recursion, we compute $Y_i \in GL_{14-i}(D)$, $i = 2, \dots, 5$ so that we finally get:

$$\underbrace{\left(\begin{array}{cc|ccc} 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ \hline 0 & 0 & \partial_1 & 0 & 0 \\ 0 & 0 & \partial_2 & -1 & 0 \\ 0 & 0 & 0 & \partial_1 & 0 \\ 0 & 0 & 0 & 1 & \partial_1 \\ 0 & 0 & 0 & 0 & \partial_2 \\ 0 & 0 & 0 & \partial_2 & 0 \end{array} \right)}_{L'_5} = Y_5^{-1} \underbrace{\left(\begin{array}{cc|ccc} \partial_1 & 0 & 0 & 0 & 0 \\ \frac{1}{2} \partial_2 & \frac{1}{2} \partial_1 & 0 & 0 & 0 \\ 0 & \partial_2 & 0 & 0 & 0 \\ \hline 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ \hline 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{array} \right)}_{L_5} X.$$

Procedure to remove identity blocks: general description

◇ Hyp.: we have computed $Y_{s,j-1} \in \text{GL}_{n-s-j+1}(D)$, $X_{j-1} \in \text{GL}_{m-j+1}(D)$ satisfying $L'_{s,j-1} = Y_{s,j-1}^{-1} L_{s,j-1} X_{j-1}$.

$$\begin{array}{ccc}
 D^{1 \times (n-s-j)} & \xrightarrow{\cdot L'_{s,j}} & D^{1 \times (m-j)} \\
 \uparrow \cdot G'_{1,j} & & \uparrow \cdot G'_{2,j} \\
 D^{1 \times (n-s-j+1)} & \xrightarrow{\cdot L'_{s,j-1}} & D^{1 \times (m-j+1)} \\
 \uparrow \cdot Y_{s,j-1} & & \uparrow \cdot X_{j-1} \\
 D^{1 \times (n-s-j+1)} & \xrightarrow{\cdot L_{s,j-1}} & D^{1 \times (m-j+1)} \\
 \uparrow \cdot W_{1,j} & & \uparrow \cdot W_{2,j} \\
 D^{1 \times (n-s-j+1)} & \xrightarrow{\cdot L_{s,j-1}} & D^{1 \times (m-j+1)} \\
 \uparrow \cdot H_{1,j} & & \uparrow \cdot H_{2,j} \\
 D^{1 \times (n-s-j)} & \xrightarrow{\cdot L_{s,j}} & D^{1 \times (m-j)}
 \end{array}$$

$$\implies L_{s,j} \underbrace{(H_{2,j} W_{2,j} X_{j-1} G'_{2,j})}_{X_j \in \text{GL}_{m-j}(D)} = \underbrace{(H_{1,j} W_{1,j} Y_{s,j-1} G'_{1,j})}_{Y_{s,j} \in \text{GL}_{n-s-j}(D)} L'_{s,j}.$$

Procedure to remove identity blocks: matrix computations

◇ The process is the same as for removing zero rows:

- 1 We decompose $Y_{s,j-1}$, X_{j-1} , and their inverses by blocks;
- 2 The key lemma and the assumption $\text{sr}(D) \leq p' - r$ implies that **there exist $c \in D$, $u \in D^{p'-j}$, and $v \in D^{1 \times (p'-j)}$** s.t., if $k_2 = p - j + 1$,

$$\begin{pmatrix} c(X'_{11})_{k_2} & v \end{pmatrix} \begin{pmatrix} (X_{11})_{\cdot k_2} \\ (X_{21})_{\cdot k_2} + u(X_{31})_{k_2} \end{pmatrix} = 1;$$

- 3 Explicit formulas for all the matrices are then obtained in terms of a solution $c \in D$, $u \in D^{p'-j}$, and $v \in D^{1 \times (p'-j)}$.

Example in the theory of 2D linear elasticity

- ◇ We have $q = 3$, $p = 2$, $q' = 6$, and $p' = 3$ so that $n = 14$ and $m = 5$.
- ◇ We have $\text{sr}(\mathbb{Q}[\partial_1, \partial_2]) = 3$.
- ◇ The positive integer r has to satisfy:

$$r \leq \min(p, p') = 2, \quad \text{sr}(D) \leq \max(p - r, p' - r) = 3 - r,$$

⇒ No positive integer r satisfies the hypothesis of Warfield's theorem.

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- ◇ But, in the above process, **the condition on r is just a sufficient condition for $(c(X'_{11})_{k_2}, v) \left(\begin{matrix} (X_{11})_{.k_2} \\ (X_{21})_{.k_2} + u(X_{31})_{k_2} \end{matrix} \right) = 1$ to admit a solution!**

⇒ In some cases, such c , u , and v could exist without the hypothesis on r .

Example in the theory of 2D linear elasticity

◇ Here, to remove the first identity block, we are reduced to solving:

$$(c \ (0 \ 0) \ v) \begin{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} \\ \begin{pmatrix} 0 \\ 0 \end{pmatrix} - u \end{pmatrix} = 1.$$

◇ Even if the hypothesis of Warfield's theorem is not fulfilled, a solution is clearly given by $c = 0$, $v = (0 \ -1)$, and $u = (0 \ 1)^T$.

◇ This allows to remove a first identity block!

Example in the theory of 2D linear elasticity

- ◇ From the explicit formulas, we get $X_1 \in GL_4(D)$ and $Y_{5,1} \in GL_8(D)$ so that we have the equivalence of matrices $L'_{5,1} = Y_{5,1}^{-1} L'_{5,1} X_1$.
- ◇ Similarly, we can remove a second identity block and we finally get:

$$\underbrace{\begin{pmatrix} 0 & 0 & 0 \\ \partial_1 & 0 & 0 \\ \partial_2 & -1 & 0 \\ 0 & \partial_1 & 0 \\ 0 & 1 & \partial_1 \\ 0 & 0 & \partial_2 \\ 0 & \partial_2 & 0 \end{pmatrix}}_{L'_{5,2}} = Y_{5,2}^{-1} \underbrace{\begin{pmatrix} \partial_1 & 0 & | & 0 \\ \frac{1}{2} \partial_2 & \frac{1}{2} \partial_1 & | & 0 \\ 0 & \partial_2 & | & 0 \\ \hline 0 & 0 & | & 1 \\ 0 & 0 & | & 0 \\ 0 & 0 & | & 0 \\ 0 & 0 & | & 0 \end{pmatrix}}_{L_{5,2}} X_2.$$

Summary

- ◇ We have given **constructive versions of Fitting and Warfield's theorems**.
- ◇ **Explicit formulas** for all unimodular matrices providing the equivalences are given in terms of the matrices defining the D -module isomorphism.
- ◇ Concerning Warfield's theorem, the method relies on the **resolution of a "stable rank" equation**.
- ◇ We have an **implementation** of all the algorithms **in Maple**. It uses heuristics for solving the "stable rank" equations which allow to treat many examples.

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Thank you for your attention!